

# 國立中央大學八十九學年度轉學生入學試題卷

資訊工程學系 三年級 科目：資料結構 共 2 頁 第 1 頁

1. Given the following definitions:

```
#define MAX_STACK 100
typedef struct stack_type {
    ITEM_TYPE item [MAX_STACK];
    int top;
} STACK_TYPE;
```

The definition of ITEM\_TYPE is unspecified and left to the stack user.

Please implement the stack functions below. (30%)

```
void create_stack(STACK_TYPE stack); /* Make stack logically accessible*/
void destroy_stack(STACK_TYPE *stack); /* Make stack logically accessible*/
BOOLEAN empty_stack(STACK_TYPE stack); /* True if stack is empty */
BOOLEAN full_stack(STACK_TYPE *stack); /* True if stack is full */
void push (STACK_TYPE stack, ITEM_TYPE new_item);
/* Add item to the top of the stack */
void pop (STACK_TYPE *stack, ITEM_TYPE *old_item);
/* Remove item from the top of the stack */
```

2. Since precedence plays an important role in transforming infix to postfix, let us assume the existence of a function pred(op1, op2), where op1 and op2 are characters representing operators. This function returns TRUE if op1 has precedence over op2 when op1 appears to the left of op2 in an infix expression without parentheses.

pred(op1, op2) returns FALSE otherwise. For example, pred('\*', '+') and pred('+', '+') are TRUE, whereas pred('+', '\*') is FALSE. To use the function to accommodate parentheses, please set the following precedence rules for parentheses using TRUE or FALSE: (20%)

pred('(', op) =	for any operator op
pred(op, '(') =	for any operator op other than ')'
pred(op, ')') =	for any operator op other than '('
pred(')', op) =	for any operator op

3. Consider a data structure to represent the queue. A queue node consists of an information field and a field holding a pointer to the next node. Given the following declarations:

```
typedef struct node_type {
    ITEM_TYPE info;
    struct node_type *next;
} NODE_TYPE;
typedef struct {
    NODE_TYPE *front;
    NODE_TYPE *rear;
} Q_TYPE;
```

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Please fill the following blanks in the implementation of the operation enqueue. The empty\_queue (queue) returns true if the queue is empty. (25%)

void enqueue (Q\_TYPE \*queue, ITEM\_TYPE item) /add a new item to the rear of the queue \*/

{

    NODE\_PTR new\_node;

    new\_node = ( NODE\_PTR ) malloc (sizeof (NODE\_TYPE));

    if (new\_node != NULL) {

        \_\_\_\_\_(1)\_\_\_\_\_;

        \_\_\_\_\_(2)\_\_\_\_\_;

        if (empty\_queue (queue) == TRUE)

            \_\_\_\_\_(3)\_\_\_\_\_;

        else

            \_\_\_\_\_(4)\_\_\_\_\_;

            \_\_\_\_\_(5)\_\_\_\_\_;

}

}

4. Given a strictly binary tree  $t$  in which the  $n$  leaves are labeled as nodes 1 through  $n$ , let  $level(i)$  be the level of node  $i$  and let  $freq(i)$  be an integer assigned to node  $i$ . Define the weighted path length of  $t$  as the sum of  $freq(i) * level(i)$  over all leaves of  $t$ . Which one of the following is the strictly binary tree with minimum weighted path length. (a) Huffman tree (b) Binary search tree (c) Heap tree (d) Threaded binary tree (5%)

5. We are given a set of 6 positive weights 2, 3, 5, 7, 9, and 13. Exactly one of these weights is to be associated with each of the 6 external nodes in a binary tree with 5 internal nodes. The weighted external path length of such a binary tree is defined to

$$\sum_{i=1}^6 q_i k_i$$
 where  $k_i$  is the distance from the root node to the external node with weight  $q_i$ . Please compute the minimal weighted external path length of the tree. (20%)